

(Im-)Proving Landauer's Principle

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arXiv:1306.xxxx, Landauer's Principle

arXiv:1304.0036, Lower bound on relative entropy
(and applications)

Landauer's Principle – a common formulation

Suppose a computer “erases” 1 bit of information.

Then: The amount of “heat” “dissipated” into the environment is at least $k_B T \log 2$:

$$\Delta Q \geq k_B T \log 2 ,$$

where T = temperature of the environment of the computer.

$$\beta\Delta Q \geq \Delta S \quad \text{“Landauer bound”}$$

$$\text{where } k_B \equiv 1, \beta \equiv 1/T$$

Why “erasure”? E.g. to re-initialize error correcting mechanism.

Existing derivations of LP

- **based on 2nd Law of Thermodyn:** e.g. Landauer '61, ...
→ mix-up of notions (cf. Earman/Norton, Bennett, ...)
- **in specific models:** e.g. 1-particle gas in box
→ need to accept thermodyn formalism (e.g. “quasistatic”)
- **recently: (more) microscopic**
 - Shizume (1995): *effective* dissipative force (Fokker-Planck)
 - Piechocinska (2000): *Jarzynski equality*
 - assumes: final product state $\rho_S \otimes \rho_R \xrightarrow{U} \rho'_S \otimes \rho'_R$
 - assumes: ρ'_S pure
 - assumes: ρ'_R diagonal in energy eigenbasis → *quantum?*
 - Sagawa/Ueda (2009): need system Hamiltonian H_S , ...
- **claimed “violations” of LP:**
→ Nieuwenhuizen '01, Lutz '11, Orlov '12, ...

our work: rigorous and minimal formulation & proof of LP

1 Formulation & Proof

- minimal setup
- LP equality: $\beta\Delta Q = \Delta S + I(S' : R') + D(\rho'_R \| \rho_R)$

2 Finite-size effects

- entropy inequalities
- finite-size effects: $\beta\Delta Q \geq \Delta S + \frac{(\Delta S)^2}{\log^2 d}, \dots$

3 Sharpness of $\beta\Delta Q \geq \Delta S$

Minimal assumptions for LP

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(1) system S , reservoir R : $\mathcal{H}_{SR} = \mathcal{H}_S \otimes \mathcal{H}_R$

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- (2) initial state uncorrelated: $\rho_{SR} = \rho_S \otimes \rho_R$

“Counterexample”: perfect classical correlations

Suppose: $\rho_{SR} = \sum_i p_i |i\rangle_S \langle i| \otimes |i\rangle_R \langle i|$.

Let: $U: |i\rangle_S |i\rangle_R \mapsto |0\rangle_S |i\rangle_R$.

Then: $U\rho_{SR}U^\dagger = |0\rangle_S\langle 0| \otimes \sum_i p_i |i\rangle_R\langle i|$

$$\rightarrow \rho'_R = \rho_R$$

→ no heat change

→ LP “violated”

Minimal assumptions for LP

- (1) system S , reservoir R : $\mathcal{H}_{SR} = \mathcal{H}_S \otimes \mathcal{H}_R$
- (2) initial state uncorrelated: $\rho_{SR} = \rho_S \otimes \rho_R$
- (3) $\rho_R = \frac{e^{-\beta H}}{\text{tr}[e^{-\beta H}]}$ (R -Hamiltonian H , R -temperature $T \equiv 1/\beta$)
- (4) unitary evolution: $\rho'_{SR} = U \rho_{SR} U^\dagger$

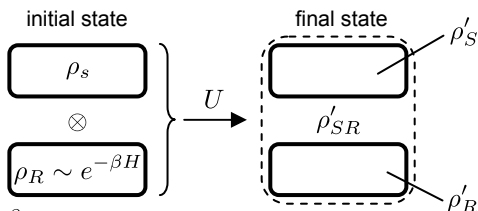
- microscopic laws of nature
- avoid obscuring entropy sinks

Minimal assumptions for LP

- (1) system S , reservoir R : $\mathcal{H}_{SR} = \mathcal{H}_S \otimes \mathcal{H}_R$
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- (4) unitary evolution: $\rho'_{SR} = U\rho_{SR}U^\dagger$

-
- no Hamiltonian for S
 - no temperature for S
 - ρ'_{SR} may be correlated
 - ρ'_S need not be pure
 - entropy of S may decrease or increase
 - quantum ($[H, \rho'_R] \neq 0$) and classical

Minimal setup for LP



- $\Delta S := S(\rho_S) - S(\rho'_S)$
- $\Delta := S(\rho'_R) - S(\rho_R)$
- $\Delta Q := \text{Tr}[H\rho'_R] - \text{Tr}[H\rho_R]$

[von Neumann entropy: $S(\rho) := -\text{Tr}[\rho \log \rho] = -\sum_i p_i \log p_i$

\rightarrow averaged quantities; also: $\Delta Q = \text{averaged heat flow}$]

Proof of LP

① “Second Law Lemma”: $\Delta = \Delta S + I(S' : R') \geq \Delta S$

Proof

$$\begin{aligned}\Delta - \Delta S &= S(\rho'_R) - S(\rho_R) + S(\rho'_S) - S(\rho_S) \\ &= S(\rho'_S) + S(\rho'_R) - S(\rho_{SR}) \\ &= S(\rho'_S) + S(\rho'_R) - S(\rho'_{SR}) \\ &= I(S' : R') \geq 0.\end{aligned}$$

- $I(A : B) \geq 0$ mutual information
- no thermal state assumption

Proof of LP

- 1 “Second Law Lemma”: $\Delta = \Delta S + I(S' : R') \geq \Delta S$
- 2 R -entropy vs. heat: $\beta\Delta Q = \Delta + D(\rho'_R \| \rho_R)$

Proof

$$\begin{aligned}\Delta &= S(\rho'_R) - S(\rho_R) \\ &= \text{tr} \left[-\rho'_R \log \rho'_R + \rho_R \log \frac{e^{-\beta H}}{\text{Tr}[e^{-\beta H}]} \right] \\ &= \text{tr} [\beta H(\rho'_R - \rho_R)] + \text{tr} \left[\rho'_R \log \frac{e^{-\beta H}}{\text{Tr}[e^{-\beta H}]} - \rho'_R \log \rho'_R \right] \\ &= \beta\Delta Q - D(\rho'_R \| \rho_R) .\end{aligned}$$

recall: $D(\sigma \| \rho) := \text{Tr}[\sigma \log \sigma] - \text{Tr}[\sigma \log \rho] \geq 0$
 (“relative entropy”)

Main result I: Equality form of LP

Theorem: Equality form of Landauer's Principle

Let $\rho_{SR} = \rho_S \otimes \rho_R$ be a product state,

where $\rho_R = e^{-\beta H} / \text{Tr} [e^{-\beta H}]$ is thermal state of Hamiltonian H at inverse temperature β .

Assume $\rho'_{SR} := U\rho_{SR}U^\dagger$ with a unitary evolution U .

Then:

$$\begin{aligned}\beta\Delta Q &= \Delta S + I(S' : R') + D(\rho'_{SR} \| \rho_{SR}) \\ &\geq \Delta S.\end{aligned}$$

Equality cases in Landauer's bound: $\beta\Delta Q = \Delta S$

Equality form of LP: $\beta\Delta Q = \Delta S + I(S' : R') + D(\rho'_R \| \rho_R)$

- $D(\rho'_R \| \rho_R) = 0 \Rightarrow \rho'_R = \rho_R$
- $I(S' : R') = 0 \Rightarrow \rho'_{SR} = \rho'_S \otimes \rho_R = U(\rho_S \otimes \rho_R)U^\dagger$

Thus: $\beta\Delta Q = \Delta S \Leftrightarrow \rho'_S = V\rho_S V^\dagger$ and $\rho'_R = \rho_R$
 $\Leftrightarrow \Delta S = \Delta Q = 0$
 \Leftrightarrow basically nothing happens

Next: explicit improvements of Landauer's bound:

→ need: finite size $d = \dim(R) < \infty$

Finite-size effects

Reservoir: $\dim(\mathcal{H}_R) = d < \infty$:

- e.g. when error-correcting mechanism small
- e.g. when short interaction time $S - R$: effectively small d

Idea: $\beta\Delta Q \geq \Delta S + D(\rho'_R \parallel \rho_R)$

$$\Delta S > 0: \Rightarrow 0 < \Delta = S(\rho'_R) - S(\rho_R)$$

$$\Rightarrow \rho'_R \neq \rho_R$$

$$\Rightarrow D(\rho'_R \parallel \rho_R) > 0.$$

→ new entropy inequality: $D(\rho'_R \parallel \rho_R) \geq M(\Delta, d)$

$$\Delta S < 0: D(\rho'_R \parallel \rho_R) \geq \frac{(\beta\Delta Q)^2}{2\max_T C(T)} \geq \frac{(\beta\Delta Q)^2}{2N(d)}$$

Relative entropy vs. entropy difference

Theorem

Let σ, ρ on \mathbb{C}^d . Define $\Delta := S(\sigma) - S(\rho)$. **Then:**

$$D(\sigma \parallel \rho) \geq M(\Delta, d) \geq \frac{\Delta^2}{2N} + \frac{\Delta^3}{6N^2} \geq 0,$$

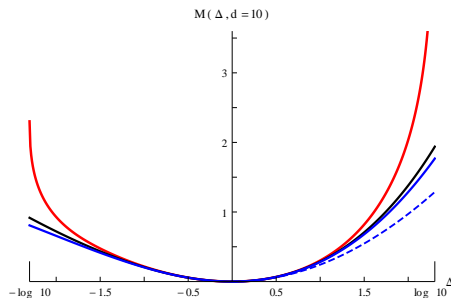
where

$$M(\Delta, d) := \min_{s,r} \left\{ D_2(s \parallel r) \mid H_2(s) - H_2(r) + (s - r) \log(d - 1) = \Delta \right\}$$

$$\text{and } N = \frac{1}{4} \log^2(d - 1) + 1 \quad \text{or} \quad N = \log^2 d.$$

- $M(\Delta, d)$: tight, effectively computable, strictly convex
- $N(d) := \max_{0 < r < 1/2} r(1 - r) \left(\log \frac{1-r}{r} (d - 1) \right)^2$: cubic Taylor

Relative entropy vs. entropy difference



graph for
 $d = 10 = \dim(R)$

$$\begin{aligned}
 D(\rho'_R \| \rho_R) &\geq M(\Delta, d) \\
 &\geq \frac{\Delta^2}{2N(d)} + \frac{\Delta^3}{6N(d)^2} \\
 &\geq \frac{\Delta^2}{\frac{1}{2} \log^2(d-1) + 2} \\
 &\geq 0.
 \end{aligned}$$

(better than Pinsker +
Fannes–Audenaert)

Relative entropy vs. entropy difference

Proof

Let $\Delta = S(\sigma) - S(\rho)$ be given:

- ① $D(\sigma \parallel \rho) = \text{Tr} [(-\log \rho) \sigma] - S(\sigma)$ for fixed ρ , fixed $S(\sigma)$
 \Rightarrow same exponential family: $\sigma = \rho^\gamma / \text{Tr} [\rho^\gamma]$
- ② Lagrange multipliers $\Rightarrow \rho, \sigma$ have at most 2 distinct EVs $\neq 0$
- ③ discrete optimization \Rightarrow 1 large EV, $(d-1)$ small EVs

$$\sigma = \text{diag} \left(1-s, \frac{s}{d-1}, \dots, \frac{s}{d-1} \right)$$

$$\rho = \text{diag} \left(1-r, \frac{r}{d-1}, \dots, \frac{r}{d-1} \right)$$

classical states
(commuting)

$$\Rightarrow M(\Delta, d) = \inf_{0 \leq s, r \leq 1} \{ D_2(s \parallel r) \mid H_2(s) - H_2(r) + (s-r) \log(d-1) = \Delta \}$$

Improved Landauer bound for $\Delta S \geq 0$

$$\begin{aligned}\beta\Delta Q &\geq \Delta S + D(\rho'_R \| \rho_R) \geq \Delta S + M(\Delta S, d) \\ &\geq \Delta S + \frac{(\Delta S)^2}{\frac{1}{2} \log^2(d-1) + 2}.\end{aligned}$$

Example: Erasure of 1 bit of information

- 2-qubit reservoir: > 34% more heat dissipation
 - 5-qubit reservoir: > 12% more heat dissipation
-
- bound tight
 - denominator $\sim \log^2 d \sim (\#particles)^2$
 - $\log d$ = effective reservoir D.O.F.'s in *short* interaction

Finite-size effects for $\beta\Delta Q \leq 0$

$$\begin{aligned}
D(\rho'_R \| \rho_R) &\geq D(\rho'_{R,th} \| \rho_R) = \beta\Delta Q - [S(\rho'_{R,th}) - S(\rho_R)] \\
&= \dots \geq \\
&\geq \int_{E_R}^{E_R + \Delta Q} \int_{E_R}^E \frac{\beta^2}{C_H(E')} dE' dE \\
&\geq \frac{(\beta\Delta Q)^2}{2N(d)}
\end{aligned}$$

[next slide: $C_H(E') \leq N(d) \leq \frac{1}{4} \log^2(d-1) + 1$]

$$\Rightarrow \beta\Delta Q \geq \Delta S + \frac{(\beta\Delta Q)^2}{2N(d)}$$

$$\Rightarrow \beta\Delta Q \geq \Delta S + \left[N(d) - \Delta S + \sqrt{N(d)^2 - 2N(d)\Delta S} \right]$$

Maximum heat capacity in $d < \infty$ dimensions

$$\begin{aligned}
 C_H(\beta) &= \frac{d}{dT} \operatorname{tr} \left[H \frac{e^{-H/T}}{\operatorname{tr}[e^{-H/T}]} \right] \bigg|_{T=\frac{1}{\beta}} = \operatorname{var}_{\rho_\beta}(\beta H) = \operatorname{var}_{\rho_\beta}(\log \rho_\beta) \\
 &= \operatorname{Tr} \left[\rho_\beta (\log \rho_\beta + S(\rho_\beta) \mathbb{1})^2 \right]
 \end{aligned}$$

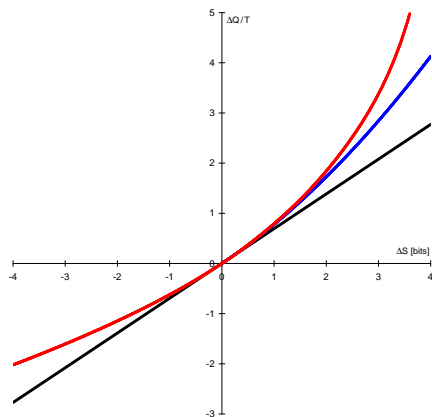
Theorem: For any state ρ on \mathbb{C}^d :

$$\begin{aligned}
 \operatorname{var}_\rho(\log \rho) &\leq N(d) := \max_{0 \leq r \leq 1/2} r(1-r) \left(\log \frac{1-r}{r} (d-1) \right)^2 \\
 &< \frac{1}{4} \log^2(d-1) + 1 \lesssim n^2 \quad \text{"superextensive"}
 \end{aligned}$$

attained for:

$$\begin{aligned}
 \rho &= \operatorname{diag} \left(1-r, \frac{r}{d-1}, \dots, \frac{r}{d-1} \right) \\
 H &= \operatorname{diag}(-1, 0, \dots, 0)
 \end{aligned}$$

Main result II: Finite-size improvements of LP



$d = 16 = \dim(R)$ (4-qubit reservoir)

- Landauer's bound:
 $\beta\Delta Q \geq \Delta S$
- tight bound for $\Delta S \geq 0$:
 $\beta\Delta Q \geq \Delta S + M(\Delta S, d)$
- for $\Delta S \leq 0$ (not tight):
 $\beta\Delta Q \geq \Delta S + \left[N_d - \Delta S - \sqrt{N_d^2 - 2N_d\Delta S} \right]$
- quadratic bound $\Delta S \geq 0$:
 $\beta\Delta Q \geq \Delta S + \frac{2(\Delta S)^2}{\log^2(d-1)+4}$

Extension: Processes with memory

Let: $\rho_{SM} \otimes \rho_R \mapsto \rho'_{SMR} = U(\rho_{SM} \otimes \rho_R)U^\dagger$

Example: perfect classical correlations

$$\rho_{SM} = \sum_i p_i |i\rangle_S \langle i| \otimes |i\rangle_M \langle i| \xrightarrow{U_{SM}} |0\rangle_S \langle 0| \otimes \sum_i p_i |i\rangle_M \langle i|$$

whereas $\rho'_R = \rho_R$, i.e. $\Delta = \Delta Q = 0$.

Result: If $\rho'_M = \rho_M$ [or $S(\rho'_M) \leq S(\rho_M)$], then:

- ① “2nd Law”: $\Delta \geq \Delta S_{cond} := S(S|M) - S(S'|M')$
- ② LP: $\beta\Delta Q \geq \Delta S_{cond} + I(S'M' : R') + D(\rho'_R || \rho_R) \geq \Delta S_{cond}$
→ **proofs** similar as before
- ③ **finite-size corrections** with ΔS_{cond} rather than ΔS

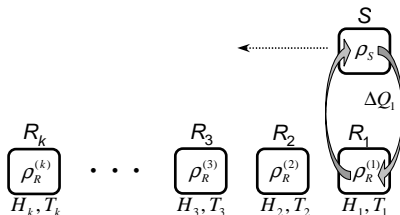
Achievability of Landauer's bound

- How sharp can $\beta\Delta Q \geq \Delta S$ be?
- Given ρ_S and ρ'_S :
Construct process with $\beta\Delta Q \rightarrow \Delta S \equiv S(\rho_S) - S(\rho'_S)$!
 \Rightarrow need $d \rightarrow \infty$

Explicit process next:

- iterated SWAP processes (for given ρ_S, ρ'_S)

Iterated SWAP process



- $\Delta S_i = S(\rho^{(i-1)}) - S(\rho^{(i)})$
- $\beta\Delta Q_i = \Delta S_i + D(\rho^{(i-1)}\|\rho^{(i)})$
- $\Delta S = S(\rho^{(0)}) - S(\rho^{(k)}) = S(\rho_S) - S(\rho'_S)$

$$\beta\Delta Q = \Delta S + \sum_{i=1}^k D(\rho^{(i-1)}\|\rho^{(i)}) \gtrsim \Delta S + \frac{(\Delta S)^2}{k \cdot \log^2 d}$$

energy–time tradeoff ? attainable ? thdyn. “reversibility” ?

Iterated SWAP process

$$\sum_{i=1}^k D(\rho^{(i-1)} \parallel \rho^{(i)}) = - \sum_{i=1}^k \text{tr}[\rho^{(i-1)} (\log \rho^{(i)} - \log \rho^{(i-1)})]$$

$$\xrightarrow{(k \rightarrow \infty)} - \sum_{i=1}^k \text{tr}[\rho^{(i-1)} ((\rho^{(i-1)})^{-\frac{1}{2}} \delta \rho_i (\rho^{(i-1)})^{-\frac{1}{2}})] = 0.$$

\Rightarrow Landauer bound **sharp** with $d = d_S^k \rightarrow \infty$

- Anders/Giovannetti (2012): $\rho^{(i)} := \frac{i}{k} \rho'_S + \frac{k-i}{k} \rho_S$. Then:

$$\sum_{i=1}^k D(\rho^{(i-1)} \parallel \rho^{(i)}) \leq \frac{D(\rho_S \parallel \rho'_S) + D(\rho'_S \parallel \rho_S)}{k}$$

\rightarrow matches our lower bound $\gtrsim (\Delta S)^2/k$ for k swaps

\rightarrow bound for general processes: $\geq M(\Delta S, d_S^k) \gtrsim (\Delta S)^2/k^2$

Caveat: What if $\text{rank}(\rho'_S) < \text{rank}(\rho_S)$?

Conclusion & technical questions

- minimal assumptions: $\rho_S \otimes e^{-\beta H} \xrightarrow{U} \rho'_{SR}$
- LP equality: $\beta\Delta Q = \Delta S + I(S' : R') + D(\rho'_R \| \rho_R)$
- finite-size effects: $\beta\Delta Q \geq \Delta S + \frac{(\Delta S)^2}{2 \log^2 d}$
 - 10% – 50% for reservoir of $N \leq 5$ qubits
 - model for energy-time tradeoff

-
-
- tight bound for $\Delta S < 0$? [possibly $\Delta S < -\log d$]
 - take $I(S' : R')$ -term into account
 - $\beta\Delta Q \geq F(\rho_S, \rho'_S, d)$
 - formulation & proof for C^* -dynamical reservoir/system

Thermodynamics of information processing

von Neumann (1949): computers, biological systems, brain:
 $\gtrsim k_B T$ heat for *every* computing operation

Landauer (1961):

- heat dissipation for *logically irreversible* operations, e.g. erasure
- $k_B T \log 2$ per bit erased
- justification by 2nd Law

Bennett (1973):

- reversible computation (esp. *quantum*)
- but prone to error (\rightarrow error correction)
- energy expense for error correction
 \rightarrow year 2000: $\sim 500 k_B T$ per bit

Maxwell (1871):

Maxwell's Demon

Szilard (1929):

Szilard engine

Why erasure?

- computation result / error syndrome “0” / “1” in register S
- “unknown” to outside: $\rho_S = \begin{pmatrix} 1/2 & 0 \\ 0 & 1/2 \end{pmatrix}$, $S(\rho_S) = 1 \text{ bit}$
- before next computation / error correction:

$$\text{RESET: } \rho_S = \begin{pmatrix} 1/2 & 0 \\ 0 & 1/2 \end{pmatrix} \mapsto \rho'_S = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

- cannot do it reversibly:

$$\begin{array}{lcl} |0\rangle_S |0\rangle_M & \mapsto & |0\rangle_S |0\rangle_M \mapsto |0\rangle_S |0\rangle_M \\ |1\rangle_S |0\rangle_M & \mapsto & |1\rangle_S |1\rangle_M \mapsto |0\rangle_S |0\rangle_M ??? \end{array}$$

another way to see this: unitaries preserve spectrum

- \rightarrow need resource: e.g. reservoir at temperature T
- more generally: “replace” ρ_S by any given ρ'_S

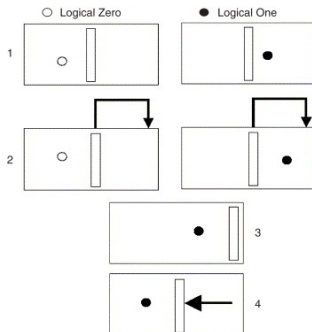
Derivation: using 2nd Law (Landauer, ...)

- ① want: possible states “0” / “1” \mapsto definite final state “0”
- ② $\Rightarrow S_{\text{system}}$ decreases by $\log 2$
- ③ \Rightarrow by 2nd Law: $S_{\text{environ}}^{\text{final}} - S_{\text{environ}}^{\text{initial}} \geq \log 2$
- ④ \Rightarrow by Th-dynamics: $\Delta Q_{\text{environ}} \geq k_B T \log 2$.

more generally: **all logically irreversible operations**

- bit erasure: $a \mapsto 0$
- AND: $(a, b) \mapsto (a, a \text{ AND } b)$
- OR: $(a, b) \mapsto (a, a \text{ OR } b)$
- ...

Derivation: 1-particle gas



3 → 4: isothermal compression

$$\begin{aligned}
 \Delta Q &= -\Delta W \\
 &= -\int_V^{V/2} p(V') dV' \\
 &= -\int_V^{V/2} \frac{1 \cdot k_B T}{V'} dV' \\
 &= k_B T \log 2 .
 \end{aligned}$$

(if “quasi-static”!)

S = information-bearing d.o.f.

E = velocity, exact location, ...

Pureness of final state

$$\lambda_{\min}(\rho'_S) \geq \sum_{i=1}^d \lambda_i^\uparrow(\rho'_{SR}) = \sum_{i=1}^d \lambda_i^\uparrow(\rho_S \otimes \rho_R) \geq d \lambda_{\min}(\rho_S) \lambda_{\min}(\rho_R)$$

$$\lambda_{\min}(\rho_R) = \frac{e^{-\beta H_{\max}}}{\text{Tr}[e^{-\beta H}]} \geq \frac{e^{-\beta H_{\max}}}{d e^{-\beta H_{\min}}}$$

$$\frac{\lambda_{\min}(\rho'_S)}{\lambda_{\min}(\rho_S)} \geq e^{-\beta(H_{\max}-H_{\min})} \geq e^{-2\beta\|H\|}.$$

→ “To erase 1 qubit”, need:

- zero-temperature reservoir ($\beta = \infty$, i.e. $\beta \Delta Q = \infty$)

- formally: $H_{\max} = +\infty$

in this case: $d < \infty \Rightarrow \Delta Q = \infty$

$d = \infty \rightarrow$ later

Example 2: Erasure towards pure states

- Want: $\rho_S = \begin{pmatrix} s_1 & 0 \\ 0 & s_2 \end{pmatrix} \xrightarrow{\text{process}} \rho'_S = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$

- $\rho_S \otimes \rho_R \mapsto \rho'_S \otimes \rho'_R$ preserves rank
 $\Rightarrow \beta\Delta Q = \infty$ whenever $\dim(R) < \infty$

- For $d = \infty$: permutation of product eigenstates

$$\rho_R = (0, r_1, 0, r_2, 0, r_3, 0, r_4, 0, r_5, 0, \dots) \quad \text{over } \ell^2$$

$$\rho'_R = (0, s_1 r_1, 0, s_2 r_1, 0, s_1 r_2, 0, s_2 r_2, 0, s_1 r_3, \dots)$$

Choose: $r_1 = 0, r_3 = (1 - \varepsilon)s_1 r_2, r_4 = (1 - \varepsilon)s_2 r_2, \dots$

$$\Rightarrow D(\rho'_R \| \rho_R) = -\log(1 - \varepsilon)$$

\Rightarrow attain Landauer limit arbitrarily closely
but need $d = \infty$ and infinitely many $+\infty$ energy levels.