## Bell nonlocality is not sufficient for the security of standard device-independent quantum key distribution protocols

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joint work with Maria Balanzó-Juandó, Karol Łukanowski, Jan Kołodyński and Antonio Acín
Phys. Rev. Lett. 127, 050503

Teiko has a problem


## Quantum advantage

## $\Downarrow$

Non-classical phenomenon

# Quantum advantage 

## $\Uparrow$ ?

Non-classical phenomenon

## Quantum advantage

Device-independent quantum key distribution

$$
\Uparrow ?
$$

Non-classical phenomenon
Bell nonlocality

## Bell nonlocality




- Quantum set:

$$
\mathcal{Q}=\left\{p(a, b \mid x, y)=\operatorname{tr}\left[\rho\left(A_{a}^{x} \otimes B_{b}^{y}\right)\right]\right\}
$$

- Convex set

- Quantum set:

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- Convex set
- Local set:

$$
\begin{aligned}
\mathcal{L}=\{p(a, b \mid x, y) & =\int_{\Lambda} p_{A}(a \mid x, \lambda) p_{B}(b \mid y, \lambda) \mathrm{d} \mu(\lambda) \\
& \left.=\sum_{\lambda^{\prime}} p_{\Lambda^{\prime}}\left(\lambda^{\prime}\right) \delta_{a, f_{A}\left(x, \lambda^{\prime}\right)} \delta_{b, f_{B}\left(y, \lambda^{\prime}\right)}\right\}
\end{aligned}
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- Convex polytope, $\mathcal{L} \subsetneq \mathcal{Q}$

Device-independent quantum key distribution (DIQKD)

## DIQKD - idea

Key Distribution


$$
K_{A}=K_{B}
$$

$K_{A}$ and $K_{B}$ are random

## DIQKD - idea

## Quantum

Key Distribution


$$
K_{A}=K_{B}
$$

$K_{A}$ and $K_{B}$ are random

$$
\left|\psi_{-}\right\rangle=\frac{1}{\sqrt{2}}(|01\rangle-|10\rangle)
$$

## DIQKD - idea



$$
K_{A}=K_{B}
$$

$K_{A}$ and $K_{B}$ are random

$$
p_{A B}(a, b \mid x, y)=\operatorname{tr}\left[\rho\left(A_{a}^{x} \otimes B_{b}^{y}\right)\right]
$$

## DIQKD based on the CHSH inequality ${ }^{1}$


${ }^{1}$ Acín, Brunner, Gisin, Massar, Pironio, Scarani, Phys. Rev. Lett. 98, 230501

## DIQKD based on the CHSH inequality ${ }^{1}$



$$
x, a, b \in\{0,1\}, y \in\{0,1,2\}
$$

Settings 0 and 1 : certifying the setup (CHSH)

$$
\rho=\left|\psi_{-}\right\rangle\left\langle\psi_{-}\right|, \quad A_{0}^{0}=|0\rangle\langle 0|, \quad A_{1}^{0}=|1\rangle\langle 1|
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Setting 2 for Bob: $B_{0}^{2}=|0\rangle\langle 0|, B_{1}^{2}=|1\rangle\langle 1|$

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Setting 2 for Bob: $B_{0}^{2}=|0\rangle\langle 0|, B_{1}^{2}=|1\rangle\langle 1|$
$x=0$ and $y=2$ : perfect randomness, perfect correlation
${ }^{1}$ Acín, Brunner, Gisin, Massar, Pironio, Scarani, Phys. Rev. Lett. 98, 230501

## Standard DIQKD protocol



$$
\begin{gathered}
n \text { rounds } \\
a_{1}, a_{2}, \ldots, a_{n} \\
b_{1}, b_{2}, \ldots, b_{n} \\
n \rightarrow \infty
\end{gathered}
$$

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input announcement standard protocol

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input announcement standard protocol

$$
\begin{gathered}
K_{A}, K_{B} \\
\text { key rate: } r \\
\frac{1}{n} I\left(K_{A}: K_{B}\right)>r-\epsilon \\
\frac{1}{n} I\left(\left\{m_{j}^{A}\right\}_{j},\left\{m_{k}^{B}\right\}_{k}, E: K_{A}\right)<\epsilon
\end{gathered}
$$

## Eavesdropping - individual attacks


$n$ rounds

$$
\begin{gathered}
\rho_{1}, \rho_{2}, \ldots, \rho_{n} \\
\rho=\frac{1}{n} \sum_{j} \rho_{j} \\
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Eavesdropper's information:

$$
\left.\begin{array}{c}
\rho_{j}, A_{a}^{x}, B_{b}^{y} \\
x_{k_{j}}, y_{k_{j}}
\end{array}\right\} e_{k_{j}}
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$$
p_{A B E}(a, b, e \mid x, y)
$$

Key extraction: $m_{1}^{A}, \ldots, m_{s}^{A}, m_{1}^{B}, \ldots, m_{s}^{B}$

## Upper bounds



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Quantum correlations

$$
p_{A B}(a, b \mid x, y)
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Which ones are useful?


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Which ones are useful?

Standard DIQKD
$\Longrightarrow p_{A B E}(a, b, e \mid x, y)$

## Upper bounds



Quantum correlations

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p_{A B}(a, b \mid x, y)
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Which ones are useful?

Standard DIQKD
$\Longrightarrow p_{A B E}(a, b, e \mid x, y)$

Classical KD results

$$
r \leq I(A: B \downarrow E)
$$

## Bell nonlocality is necessary



$$
\begin{gathered}
p_{A B}^{\mathcal{L}}(a, b \mid x, y)= \\
\sum_{\lambda} p_{\Lambda}(\lambda) \delta_{a, f_{A}(x, \lambda)} \delta_{b, f_{B}(y, \lambda)}
\end{gathered}
$$

## Bell nonlocality is not sufficient

## Specific eavesdropping attack

Specific (large) family of nonlocal correlations

The convex combination attack



Observed correlation:

$$
p=q_{\mathcal{L}} p^{\mathcal{L}}+\left(1-q_{\mathcal{L}}\right) p^{\mathcal{N} \mathcal{L}}
$$

The convex combination attack



Observed correlation:

$$
p=q_{\mathcal{L}} p^{\mathcal{L}}+\left(1-q_{\mathcal{L}}\right) p^{\mathcal{N} \mathcal{L}}
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$$
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The convex combination attack



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& \Longrightarrow p_{A B E}(a, b, e \mid x, y)
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## The convex combination attack



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Maximising $q_{\mathcal{L}}$ : linear program

## The convex combination attack



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Maximising $q_{\mathcal{L}}$ : linear program

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r \leq I(A: B \downarrow E)
$$

## Protocols with Werner states and projective measurements



$$
p_{A B}(a, b \mid x, y)=\operatorname{tr}\left[\left(v\left|\psi_{-}\right\rangle\left\langle\psi_{-}\right|+(1-v) \frac{\mathbb{I}}{4}\right)\left(A_{a}^{x} \otimes B_{b}^{y}\right)\right]
$$



Designolle, Iommazzo, Besançon, Knebel, Gelß, Pokutta, arXiv:2302.04721 (see poster no. 5, Sébastien Designolle)

## Convex combination attack

$$
\begin{aligned}
& \rho^{\mathcal{L}}=v_{\mathcal{L}}\left|\psi_{-}\right\rangle\left\langle\psi_{-}\right|+\left(1-v_{\mathcal{L}}\right) \frac{\mathbb{I}}{4}, \rho^{\mathcal{N L}}=\left|\psi_{-}\right\rangle\left\langle\psi_{-}\right| \\
& q_{\mathcal{L}}=(1-v) /\left(1-v_{\mathcal{L}}\right)
\end{aligned}
$$



## Implications and limitations

All the commonly used protocols become insecure while still exhibiting nonlocality


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Multiple parties? (see poster no. 22, Jan Nöller)

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